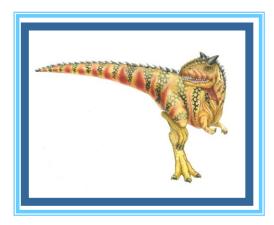
Lecture13 Bounded Buffer Problem



Modified from Silberschatz, Galvin and Gagne & Stallings

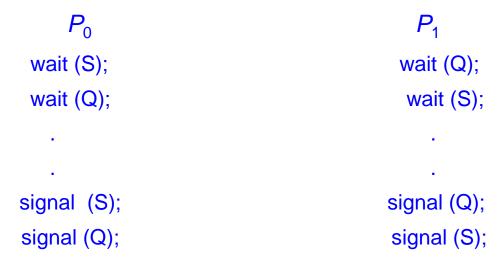


- Background
- The Critical-Section Problem
- Peterson's Solution
- Synchronization Hardware
- Semaphores
- Classic Problems of Synchronization
- Monitors
- Synchronization Examples
- Atomic Transactions



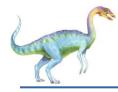


- Deadlock two or more processes are waiting indefinitely for an event that can be caused by only one of the waiting processes
- Let S and Q be two semaphores initialized to 1



Priority Inversion - Scheduling problem when lower-priority process holds a lock needed by higher-priority process





Starvation – indefinite blocking.

A process may never be removed from the semaphore queue in which it is suspended

- Order of arrival retainment:
 - Weak semaphores:
 - > The thread that will access the critical region next is selected randomly
 - Starvation is possible
 - Strong semaphores:
 - The thread that will access the critical region next is selected based on its arrival time, e.g. FIFO
 - Starvation is not possible



Other Synchronization problem(Classical Problems of Synchronization)

- Bounded-Buffer Problem
- Readers and Writers Problem
- Dining-Philosophers Problem





Bounded-Buffer Problem

- N buffers, each can hold one item
- Semaphore mutex initialized to the value 1
- Semaphore full initialized to the value 0
- Semaphore empty initialized to the value N.





The structure of the producer process

do {

// produce an item in nextp

wait (empty);
wait (mutex);

// add the item to the buffer

signal (mutex);
signal (full);
} while (TRUE);





The structure of the consumer process

do {
 wait (full);
 wait (mutex);

// remove an item from buffer to nextc

signal (mutex); signal (empty);

// consume the item in nextc

} while (TRUE);







• Q: Explain bounded buffer algorithm.

